

Aionda Mail Security Whitepaper

Zero-Knowledge · Post-Quantum · Made in Germany

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PUBLIC DOCUMENT

Contents

1. Executive Summary	3
2. Threat Model	3
3. Architecture Overview	4
4. Zero-Knowledge Authentication (OPAQUE)	5
5. Vault Master Key & Shamir Secret Sharing	7
6. Post-Quantum Hybrid KEM	8
7. Email Encryption Pipeline	9
8. Encrypted API Transport Layer	12
9. Folder Sharing Cryptography	13
10. Vault Drive — Encrypted Document Storage	14
11. Aionda Chat — Post-Quantum E2EE Messaging	20
12. Aionda Video Calls — Post-Quantum E2EE Conferencing	25
13. Side-Channel Protections	29
14. Key Management & Lifecycle	30
15. Recovery Mechanism	31
16. Passwordless Authentication (Passkeys)	32
17. Guardian: MITM Protection & Response Signing	33
18. Enterprise Email Archive (Blockchain)	36
19. What the Server Sees — and What It Does Not	40
20. Algorithm Reference	41
21. Comparison with Other Providers	42
22. Limitations & Honest Boundaries	43
23. Roadmap	44
Document History	45
Contact	46

1. Executive Summary

Aionda Mail is a zero-knowledge, post-quantum encrypted email service operated by Aionda GmbH in Stuttgart, Germany. The service combines disposable email addresses (DEAs) with a fully encrypted mailbox — a combination not offered by any other provider.

Core security properties:

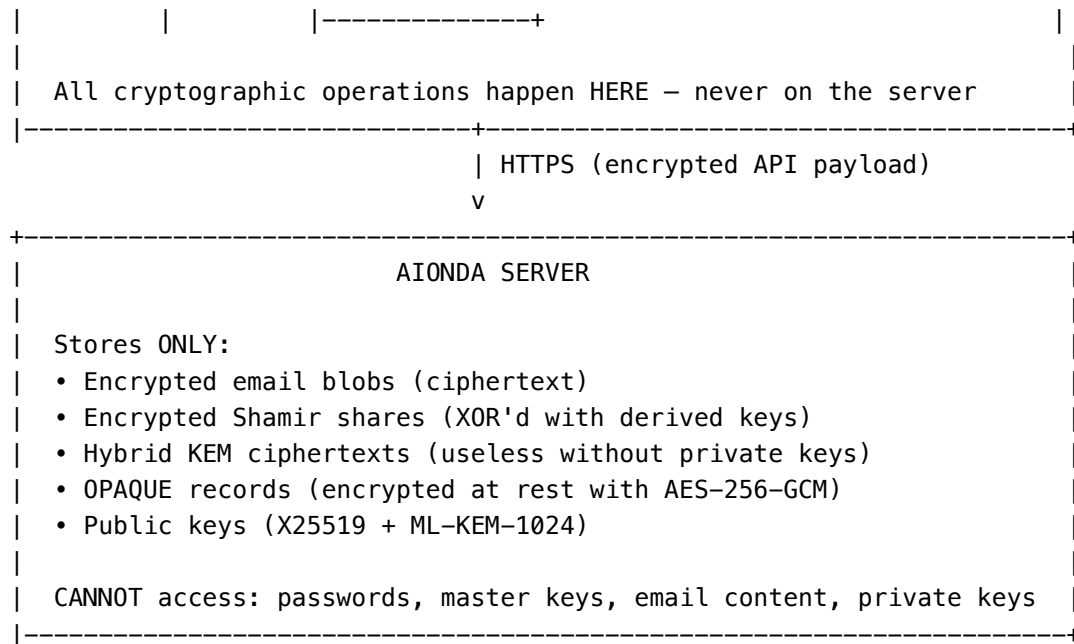
- **Zero-Knowledge Architecture:** All encryption and decryption happens exclusively in the user’s browser. The server never has access to plaintext email content in the secure mailbox, passwords, or encryption keys.
- **Post-Quantum Security:** Hybrid Key Encapsulation Mechanism (X25519 + ML-KEM-1024) protects all data against both classical and quantum computer attacks.
- **Zero-Knowledge Authentication:** OPAQUE protocol (RFC 9807) ensures passwords are never transmitted to or stored on the server — not even as hashes.
- **Shamir Secret Sharing (2-of-3):** The vault master key is split into three shares protected by password, passkey, and recovery key. Any two shares reconstruct the master key.
- **Perfect Forward Secrecy:** Every API request uses a unique, one-time cryptographic keypair. Compromising one request does not affect any other.
- **MITM Protection (Guardian):** The browser extension independently verifies all server responses via Ed25519 signatures and detects man-in-the-middle attacks through TLS certificate verification — even against corporate proxies and compromised CDNs.
- **GoBD-Compliant Email Archive:** Enterprise accounts benefit from a tamper-evident hash chain (SHA3-256 blockchain) with end-to-end encrypted content (Hybrid KEM), complete audit trail, legal hold, and configurable retention — compliant with German GoBD regulations.
- **No Password Recovery:** If the password and all recovery methods are lost, the data is irrecoverable. This is by design — it proves the server cannot access user data.

Jurisdiction: German law (DSGVO/GDPR), no data sharing with foreign intelligence agencies.

2. Threat Model

2.1 What Aionda Mail Protects Against

Threat	Protection
Server compromise (database leak, insider access)	All email content encrypted with keys the server never possesses
Network eavesdropping (ISP, Wi-Fi, CDN)	End-to-end encrypted API transport via Hybrid KEM
CloudFlare inspection	API requests are encrypted before leaving the browser; CloudFlare sees only ciphertext. Guardian extension detects response tampering via Ed25519 signatures
Corporate MITM proxies (ZScaler, Fortinet, etc.)	Guardian extension detects proxy certificates via issuer blocklist (Firefox)
Quantum computer attacks (“harvest now, decrypt later”)	ML-KEM-1024 (NIST FIPS 203) provides post-quantum resistance



3.1 Component Overview

Component	Purpose	Location
OPAQUE (RFC 9807)	Zero-knowledge password authentication	Client + Server
Shamir Secret Sharing	Vault master key protection (2-of-3 threshold)	Client only
Hybrid KEM (X25519 + ML-KEM-1024)	Post-quantum key encapsulation for emails	Client + Server
AES-256-GCM	Authenticated symmetric encryption	Client + Server
HKDF-SHA256	Key derivation from hybrid shared secrets	Client + Server
BIP39	Recovery key encoding (24-word mnemonic)	Client only
WebAuthn PRF	Passkey-based vault unlock	Client only
Bucket Padding	Side-channel protection	Client + Server

4. Zero-Knowledge Authentication (OPAQUE)

4.1 Why Not Password Hashing?

Traditional services store password hashes (bcrypt, Argon2). While better than plaintext, this approach has fundamental weaknesses:

- The server sees the password during login (even if only briefly in RAM)
- Password hashes can be brute-forced offline if the database is stolen

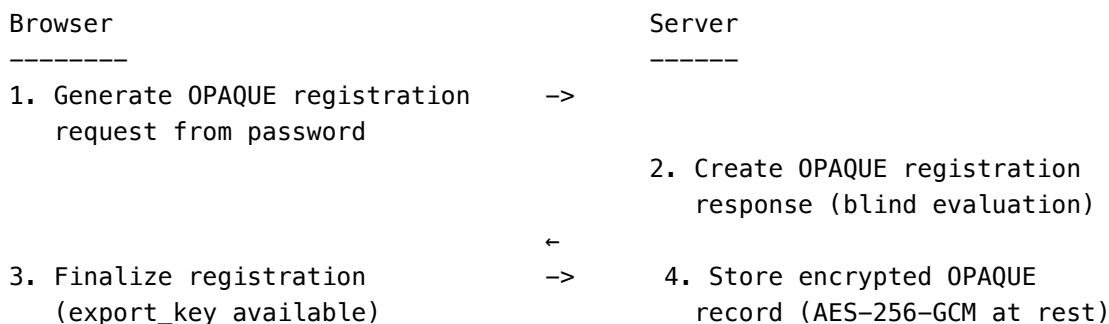
- The server could be modified to log passwords

OPAQUE eliminates all three problems. The password never leaves the browser — not as plaintext, not as a hash, not in any form.

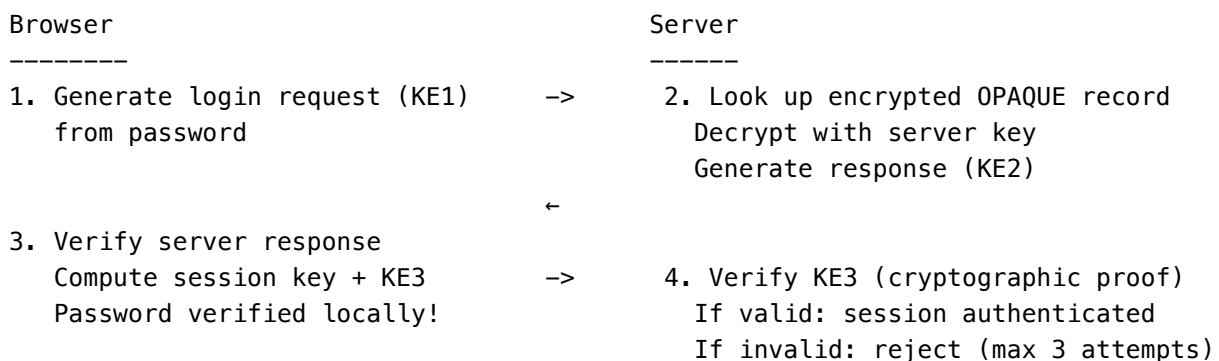
4.2 How OPAQUE Works

OPAQUE (RFC 9807) is an Asymmetric Password-Authenticated Key Exchange (aPAKE) protocol. It uses a cryptographic challenge-response mechanism where the server can verify the user knows the correct password without ever learning what that password is.

Registration (one-time):



Login (every session):



Key properties:

- Password is verified **client-side** in step 3 — the server never sees it
- The server stores an **OPAQUE record**, which is not a password hash and cannot be brute-forced offline
- OPAQUE records are additionally **encrypted at rest** with AES-256-GCM using a server-side key
- **User enumeration protection:** Non-existent accounts receive deterministic fake responses with identical timing
- **Rate limiting:** Maximum 3 authentication attempts per session, 120-second session timeout

4.3 Implementation

- **Library:** @serenity-kit/opaque (WASM-based, production-grade)
- **Server component:** Dedicated microservice for OPAQUE cryptographic operations
- **Base64 format:** base64url (URL-safe, no padding) for protocol compatibility

- **Audit logging:** All authentication events logged with timestamps and IP addresses

4.4 SRP Migration

Legacy accounts using SRP-6a are automatically migrated to OPAQUE upon next login. After migration, the SRP verifier is permanently deleted. Migration is one-way — accounts cannot revert to SRP.

5. Vault Master Key & Shamir Secret Sharing

5.1 Master Key Generation

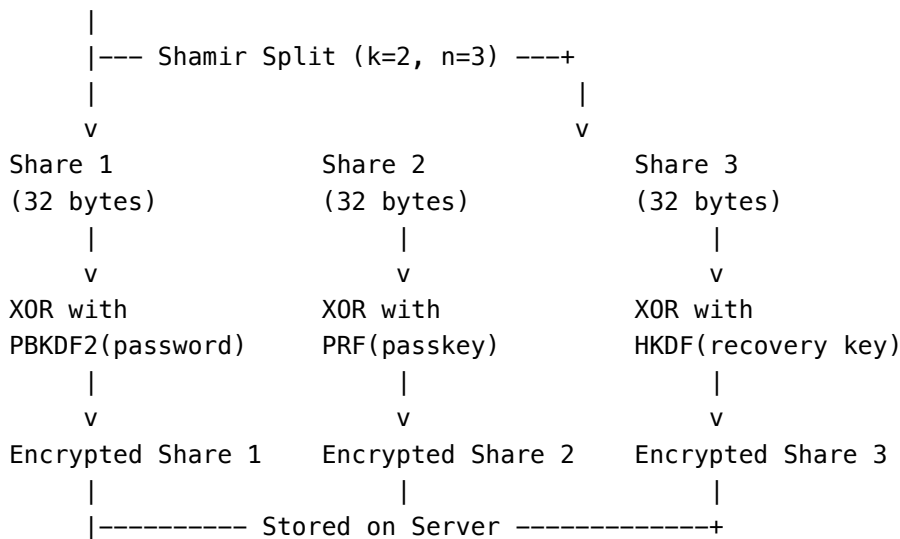
When a user activates the encrypted mailbox, a **256-bit (32-byte) master key** is generated using the browser’s cryptographically secure random number generator (`crypto.getRandomValues`).

This master key is the root of all encryption. It never leaves the browser in plaintext. It is never stored anywhere — not in the browser, not on the server, not in any form.

5.2 Shamir Secret Sharing (2-of-3)

The master key is split into three shares using **Shamir’s Secret Sharing** scheme over the Galois Field $GF(2^8)$ with the AES irreducible polynomial $(x^8 + x^4 + x^3 + x + 1)$.

Master Key (32 bytes, generated once)



Threshold property: Any 2 of the 3 shares are sufficient to reconstruct the master key via Lagrange interpolation. The server stores only the encrypted shares — and cannot decrypt any of them.

5.3 Share Protection

Each share is XOR’d with a key derived from a different authentication factor:

Share	Protected By	Key Derivation
Share 1	Password	PBKDF2-SHA256, 600,000 iterations, 32-byte random salt

Share	Protected By	Key Derivation
Share 2	Passkey (FIDO2)	WebAuthn PRF extension, hardware-bound
Share 3	Recovery Key	HKDF-SHA3-256 with account-bound salt

Reconstruction scenarios:

- **Normal login:** Password (Share 1) + Passkey (Share 2) -> Master Key
- **Lost passkey:** Password (Share 1) + Recovery Key (Share 3) -> Master Key
- **Password change:** Passkey (Share 2) + Recovery Key (Share 3) -> Master Key

5.4 Session Storage Protection

Even within a browser session, the master key is never stored in plaintext:

1. An ephemeral AES-256 key is generated
2. The master key is encrypted with this ephemeral key
3. Only the encrypted blob is placed in `sessionStorage`
4. The ephemeral key exists only in JavaScript memory (garbage-collected on tab close)

6. Post-Quantum Hybrid KEM

6.1 Why Post-Quantum?

Quantum computers running Shor's algorithm could break classical public-key cryptography (RSA, ECDH, X25519) in polynomial time. While large-scale quantum computers do not yet exist, the threat of “**harvest now, decrypt later**” attacks is real: adversaries could store encrypted data today and decrypt it when quantum computers become available.

6.2 Hybrid Approach

Aionda Mail uses a **hybrid Key Encapsulation Mechanism** combining:

- **X25519** (Curve25519 ECDH) — proven classical security, 128-bit security level
- **ML-KEM-1024** (NIST FIPS 203, formerly Kyber-1024) — post-quantum security, NIST Security Level 5

The hybrid approach provides **defense-in-depth**: the combined key is secure as long as **at least one** of the two algorithms remains unbroken.

6.3 Encapsulation Process

Sender (encrypting an email):

1. Generate ephemeral X25519 keypair
2. X25519 key agreement with recipient's public key
-> `x25519SharedSecret` (32 bytes)
3. ML-KEM-1024 encapsulation with recipient's public key
-> `mLKemSharedSecret` (32 bytes) + `mLKemCiphertext` (1568 bytes)

4. Combine secrets:


```
combinedSecret = x25519SharedSecret || mlKemSharedSecret (64 bytes)
```
5. Derive final key:


```
sharedSecret = HKDF-SHA256(
  ikm = combinedSecret,
  salt = nil,
  info = "trashmail-hybrid-kem-v1",
  length = 32
)
```
6. Use sharedSecret to wrap the email's ephemeral AES-256 key

6.4 Decapsulation Process

Recipient (decrypting an email):

1. X25519 key agreement:


```
x25519Shared = X25519(recipientPrivateKey, ephemeralPublicKey)
```
2. ML-KEM-1024 decapsulation:


```
mlKemShared = ML-KEM-1024.Decapsulate(mlKemCiphertext, recipientPrivateKey)
```
3. Combine and derive (identical to sender):


```
sharedSecret = HKDF-SHA256(x25519Shared || mlKemShared, "trashmail-hybrid-kem-v1")
```
4. Unwrap email's ephemeral AES-256 key using sharedSecret
5. Decrypt email content with ephemeral key

6.5 Key Sizes

Parameter	Size	Standard
X25519 public key	32 bytes	RFC 7748
X25519 private key	32 bytes	RFC 7748
ML-KEM-1024 public key	1,568 bytes	NIST FIPS 203
ML-KEM-1024 private key	3,168 bytes	NIST FIPS 203
ML-KEM-1024 ciphertext	1,568 bytes	NIST FIPS 203
Combined shared secret	32 bytes	HKDF-SHA256 output

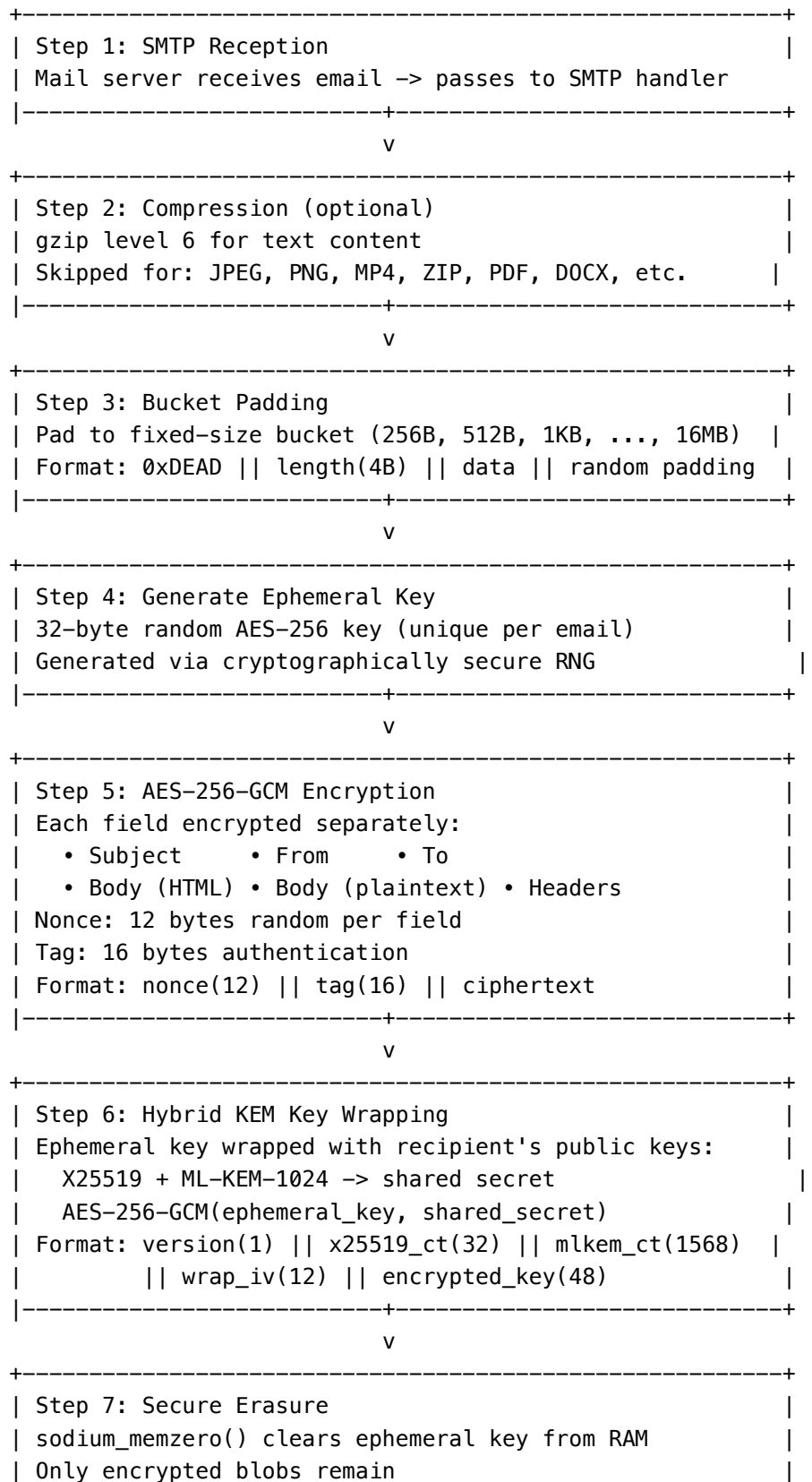
6.6 Library

- **ML-KEM-1024:** @noble/post-quantum (audited, pure JavaScript implementation)
- **X25519:** WebCrypto API (crypto.subtle.deriveBits)
- **HKDF:** @noble/hashes (RFC 5869 compliant)

7. Email Encryption Pipeline

7.1 Incoming Email (SMTP -> Encrypted Storage)

When an email arrives at Aionda Mail's SMTP server:



8. Encrypted API Transport Layer

8.1 Problem

Even with HTTPS, certain intermediaries can inspect traffic:

- **CloudFlare** (CDN/DDoS protection) terminates TLS and can see plaintext requests
- **Corporate proxies** may perform TLS inspection
- **API parameters** (like `?cmd=read_email&id=123`) leak metadata

8.2 Solution: End-to-End Encrypted API

All API communication is additionally encrypted end-to-end between the browser and the application server, inside the HTTPS tunnel:

Browser	-----	Server	-----
Phase 1: Key Exchange (once per session)			

GET /get_encryption_keys	->	Return 20 pre-generated Hybrid KEM keypairs	
	←	{uuid, x25519_pub, mlkem_pub}	
Phase 2: Encrypted Request (every API call)			

1. Pick random keypair from cache			
2. Hybrid KEM encapsulate -> shared secret			
3. gzip compress request payload			
4. Bucket-pad compressed data			
5. AES-256-GCM encrypt with shared secret			
6. Generate ephemeral response keypair			
7. POST /e {	->	8. Validate key ownership	
encrypted_payload,		9. Hybrid KEM decapsulate	
key_uuid,		10. AES-256-GCM decrypt	
x25519_ciphertext,		11. Decompress	
mlkem_ciphertext,		12. Route to API controller	
response_x25519_pub,		13. Execute business logic	
response_mlkem_pub		14. Encrypt response with client's response keys	
}	←	15. Return encrypted response	
16. Hybrid KEM decapsulate response			
17. AES-256-GCM decrypt			
18. Decompress -> plaintext response			

8.3 Key Properties

- **One-time use:** Each API keypair is used exactly once, then permanently invalidated
- **Perfect Forward Secrecy:** Compromising one request key does not affect any other request

- **Session-bound:** Keys are claimed by a specific session and cannot be reused by another
- **Key pool:** Server maintains approximately 100,000 pre-generated keypairs
- **Auto-refetch:** Client automatically requests new keys when cache drops below 10
- **Key TTL:** Claimed keys expire after 24 hours
- **Bidirectional:** Both request AND response are encrypted — the server never returns plaintext

8.4 What CloudFlare Sees

With this architecture, CloudFlare (or any TLS-terminating proxy) sees only:

- POST /e — a single, opaque endpoint
- A binary blob of encrypted data
- No API command names, no parameters, no email IDs, no user data

9. Folder Sharing Cryptography

9.1 Sharing Model

Users can share encrypted folders with other Aionda Mail users. The sharing mechanism uses the Hybrid KEM to encrypt a folder-specific key for each recipient.

9.2 Sharing Flow

Folder Owner

Recipient

- | | |
|--|---|
| <p>-----</p> <ol style="list-style-type: none"> 1. Derive folder key from master key:
 <code>folderKey = HKDF-SHA256(masterKey, folderUuid)</code> 2. Fetch recipient's public keys:
 <code>recipient.x25519_pub (32 bytes)</code>
 <code>recipient.mlkem_pub (1568 bytes)</code> 3. Hybrid KEM encapsulate:
 <code>hybridEncapsulate(recipient.x25519_pub, recipient.mlkem_pub)</code>
 <code>-> {x25519Ciphertext, mlkemCiphertext, sharedSecret}</code> 4. Encrypt folder key:
 <code>wrappedKey = AES-256-GCM(folderKey, sharedSecret, nonce)</code> 5. Store on server:
 <code>{x25519_ct, mlkem_ct, nonce, wrappedKey, permissions}</code> | <p>-----</p> <ol style="list-style-type: none"> 6. Fetch sharing record from server 7. Hybrid KEM decapsulate:
 <code>hybridDecapsulate(x25519_ct, mlkem_ct,</code>
 <code> own_x25519_priv, own_mlkem_priv)</code>
 <code>-> sharedSecret</code> |
|--|---|

8. Decrypt folder key:


```
folderKey = AES-GCM-decrypt(
    wrappedKey, sharedSecret, nonce)
```
9. Decrypt emails in folder using folderKey

9.3 Permission Model

Permission	Capability
readonly	Read folder emails (decrypt only)
einliefern	Submit new emails into folder
bearbeiten	Edit folder contents
antworten	Reply to emails within folder
vollzugriff	Full access including ownership transfer

9.4 Cross-Vault Copy (Re-Wrapping)

When a recipient copies an email from a shared folder to their own vault, the email key must be **re-wrapped** for their own Hybrid KEM keypair. This is performed entirely client-side:

1. Decrypt shared folder key using recipient's private keys
2. Decrypt email's ephemeral key using folder key
3. Re-encapsulate ephemeral key with recipient's own public keys
4. Store re-wrapped copy in recipient's vault

The server facilitates the transfer but never sees any plaintext key material.

10. Vault Drive — Encrypted Document Storage

Vault Drive is the zero-knowledge document storage of Aionda Mail. Unlike emails, which use per-message ephemeral keys, Drive uses a **per-file-key architecture**: every document receives its own 32-byte AES-256-GCM key, generated client-side and wrapped with the master key.

10.1 Why per-file keys?

The per-file-key architecture offers three key advantages:

1. **Granular sharing:** Individual documents can be shared without exposing the master key. The server simply re-wraps the file key for the recipient — the content remains unchanged.
2. **Compromise isolation:** If a single file key is compromised (e.g. via a shared link), all other documents remain protected.
3. **Efficient sharing without re-encryption:** When sharing, the content does not need to be re-encrypted — all recipients read the same ciphertext with a re-wrapped key.

10.2 Upload flow

Client

Server

- ```

```
1. Generate random file key:  
fileKey = randomBytes(32)
  2. Encrypt content, name, MIME type:  
encContent = AES-256-GCM(content, fileKey, nonce1)  
encName = AES-256-GCM(name, fileKey, nonce2)  
encMime = AES-256-GCM(mime, fileKey, nonce3)
  3. Wrap file key with master key:  
wrappedKey = AES-256-GCM(fileKey, masterKey, nonce4)
  4. Send encrypted package:  
POST /e {  
 encContent, encName, encMime,  
 wrappedKey, wrappedKeyNonce  
}
  5. Store in vault\_files + vault\_file\_chunk  
(pure ciphertexts, no key material)

### 10.3 Download flow

1. Fetch chunk + wrapped\_key from server
2. Unwrap file key:  
fileKey = AES-256-GCM-decrypt(wrappedKey, masterKey, nonce)
3. Decrypt content, name, MIME:  
content = AES-256-GCM-decrypt(encContent, fileKey, nonce1)

Decryption happens inside a **Web Worker** that holds the master key only temporarily in memory. After the operation, the memory is overwritten.

### 10.4 Sharing — key rewrap only

The central advantage of per-file keys shows when sharing: the content is **not** re-encrypted. Only the 32-byte file key is re-wrapped using the recipient's hybrid KEM.

Owner

-----

Recipient

-----

1. Unwrap file key:  
fileKey = AES-GCM-decrypt(wrappedKey, masterKey)
2. Fetch recipient's public keys:  
recipient.x25519\_pub, recipient.mlkem\_pub
3. Hybrid KEM encapsulation:  
hybridEncapsulate(recipient.x25519\_pub,  
 recipient.mlkem\_pub)  
-> {x25519\_ct, mlkem\_ct, sharedSecret}
4. Wrap file key with sharedSecret:

```
shareWrappedKey = AES-256-GCM(fileKey,
 sharedSecret, nonce)
```

5. Store share record:

```
INSERT INTO vault_file_shares {
 file_uuid, recipient_account_id,
 x25519_ephemeral, mlkem_ciphertext,
 share_wrapped_key, permissions
}
```

6. Fetch share record + chunk

7. Hybrid KEM decapsulation:

```
hybridDecapsulate(x25519_ct, mlkem_ct,
 own_x25519_priv, own_mlkem_priv)
-> sharedSecret
```

8. Unwrap file key:

```
fileKey = AES-GCM-decrypt(
 shareWrappedKey, sharedSecret)
```

9. Decrypt SAME ciphertext:

```
content = AES-GCM-decrypt(
 encContent, fileKey)
```

The owner does not need the recipient's master key — only the public hybrid KEM keys, which the server stores in plaintext.

### 10.5 Folder sharing (recursive)

When a folder is shared, the client processes all contained documents and subfolders recursively. Each file receives its own share record containing the respective document's file key, re-wrapped for the recipient. The folder itself has no folder key — the structure is linked via parent UUIDs.

**Important:** The KEM encapsulation step is performed **per operation**, not per file. All files in a batched share use the same `sharedSecret` — making the operation efficient without reducing security (each file still has its own file key).

### 10.6 Permission model

| Permission          | Can perform                                                                       |
|---------------------|-----------------------------------------------------------------------------------|
| <b>Read</b>         | Download, preview, read metadata                                                  |
| <b>Full access</b>  | + rename, upload new version, move inside the shared folder, upload new documents |
| <b>Not possible</b> | Delete (owner only), move out of shared folder (owner only)                       |

Permissions are stored in the `vault_file_shares` record and enforced server-side. Cryptography protects the content — access control protects the operations.

### 10.7 Rename, overwrite and metadata updates

For shared documents, rename and overwrite operations are performed directly on the `vault_files` record — not on individual share records. All recipients immediately see the new name or con-

tent, as they reference the same ciphertext.

On **overwrite** (new version), a **new file key** is generated, the old ciphertext is replaced, and all existing share records are re-wrapped client-side with the new key. The owner must load the public keys of all recipients for this operation.

### 10.8 Preview and in-memory decryption

Images, PDFs and other documents are decrypted **exclusively in memory** for preview. There is no disk cache with plaintext data. The client uses the **VaultDataLayer**, an encrypted IndexedDB cache that persistently stores ciphertexts and only decrypts them on demand.

Thumbnails are **never** generated server-side — the server never sees content. Thumbnails are generated client-side from the decrypted original and cached encrypted as well.

### 10.9 What the server sees

| Visible to server                       | Not visible                              |
|-----------------------------------------|------------------------------------------|
| Encrypted content (random bytes)        | Plaintext content                        |
| Encrypted filename (random bytes)       | Plaintext filename                       |
| Encrypted MIME type (random bytes)      | File type (image, PDF, text...)          |
| File size (padded to bucket boundaries) | Actual original size                     |
| Upload timestamp                        | File key, master key                     |
| Owner account ID                        | Contents of subfolders                   |
| Parent folder UUID (for structure)      | Folder names (also encrypted)            |
| Share records with KEM ciphertexts      | Recipient master key, unwrapped file key |

### 10.10 Quota enforcement

Quota enforcement happens server-side based on the **encrypted file size** (including bucket padding). The server does not know the actual plaintext size — the padding overhead is intentionally paid by the user to prevent size leakage.

Limits: - **Free**: 100 MB total storage, max. 10 MB per document - **Plus**: 1 GB total storage, max. 100 MB per document

### 10.11 External Sharing — public share links

Vault Drive documents can be shared with recipients who **do not have an Aionda Mail account** via public share links served from the isolated domain `mail.aionda.com`. The feature preserves zero-knowledge by treating each share as an **independent crypto envelope**, decoupled from the owner's vault master key.

#### Protection modes:

| Mode      | Trust factor                        | Use case                                        |
|-----------|-------------------------------------|-------------------------------------------------|
| link_only | URL fragment (never sent to server) | Convenience — anyone with the link can access   |
| password  | Argon2id-derived key                | Out-of-band password transfer (SMS, phone call) |

| Mode             | Trust factor                      | Use case                                                          |
|------------------|-----------------------------------|-------------------------------------------------------------------|
| recipient_pubkey | Hybrid KEM (X25519 + ML-KEM-1024) | Post-quantum secure when recipient has pre-registered public keys |

### 10.11.1 Share creation (owner client)

1. fileKey := AES-GCM-decrypt(vault\_files.wrapped\_key, masterKey)
2. shareKey := randomBytes(32) // ephemeral, per share
3. wrappedFileKey := AES-GCM(fileKey, shareKey) // new envelope
4. unlockVerifier := HMAC-SHA256(shareKey, "unlock:" || shareUuid) // proof-of-knowledge
5. If password protection:
  - salt := randomBytes(16)
  - pwKey := Argon2id(password, salt, t=3, m=64MB, p=1)
  - wrappedShareKey := AES-GCM(shareKey, pwKey)
  - > Link: https://mail.aionda.com/s/<shareUuid>
  - Else (link\_only):
  - > Link: https://mail.aionda.com/s/<shareUuid>#<base64(shareKey)>
  - (URL fragment – browsers never transmit fragments to the server)
6. Encrypted metadata (per-share, with shareKey):
  - encFilename, encMime, encMessage // all AES-GCM
7. POST /e (encrypted API transport, see §8)
  - > server stores share record – never sees plaintext

**Link delivery is user-chosen, not routed through Aionda Mail.** After creation, the owner decides which channel transports the link: copy-to-clipboard, QR code, a pre-filled mailto: draft opened in the user's own mail client, Signal, SMS, AirDrop, or any other out-of-band channel. Aionda Mail never sees, sends, or logs the outgoing delivery – the server only learns which share-UUIDs exist. This matters for two reasons: (a) the owner can select the most trusted channel available to them (a corporate compliance policy, a preferred messenger, a verified face-to-face QR scan), and (b) for password-protected shares it enables **channel separation** – link via channel A, password via channel B – so a single compromised channel never grants access on its own.

**10.11.2 Share access (recipient, no account required)** The recipient visits <https://mail.aionda.com/s/<shareUuid>>. The Share Page is a **separate bundle** from the Manager, with its own reproducible build, Sub-resource Integrity manifest, and a /.well-known/integrity.json endpoint for offline verification.

1. Share Page bootstrap (client generates ephemeral hybrid KEM keypair)
2. share\_fetch\_meta -> { wrapped\_file\_key, salt?, encrypted\_message, ... }
3. Unlock phase – produces a one-time unlock\_token:
  - password mode: server verifies Argon2id derivation -> unlock\_token
  - link\_only mode: client computes HMAC-SHA256(shareKey, "unlock:" || uuid)
  - server verifies against stored unlockVerifier
  - > unlock\_token
4. fileKey := AES-GCM-decrypt(wrappedFileKey, shareKey)

5. `share_download_chunk` (per chunk, `unlock_token` required)
6. Client hashes plaintext, calls `share_confirm_download`

The `unlock_token` unifies the flow across all three protection modes and enables centralized rate-limiting and audit-logging — link-only access requires proof of knowledge of the fragment, so bots and crawlers without the fragment cannot issue downloads.

**10.11.3 Enforced policies** Every share must declare the following at creation time (all enforced server-side):

- **`expires_at`** — mandatory, default 7 days, maximum 90 days
- **`max_downloads`** — mandatory counter, default 10
- **Rate limiting** — Argon2 attempts on password shares are throttled per IP hash and per share
- **`revoked_at`** — the owner can revoke any share with a single click; future unlock attempts and downloads return a unified “share unavailable” response (same error as expired or exhausted — no enumeration oracle)

**10.11.4 Tamper-evident audit chain** All relevant access events are appended to the existing `enterprise_audit_log` hash chain (SHA3-256, see §18.2). The following action types are reserved for external sharing:

`EXT_SHARE_CREATED`, `EXT_SHARE_EMAIL_SENT`, `EXT_SHARE_VIEWED`, `EXT_SHARE_UNLOCKED`, `EXT_SHARE_PREVIEWED`, `EXT_SHARE_DOWNLOAD_STARTED`, `EXT_SHARE_DOWNLOAD_CONFIRMED`, `EXT_SHARE_INTEGRITY_BROKEN`, `EXT_SHARE_PASSWORD_FAIL`, `EXT_SHARE_REVOKED`, `EXT_SHARE_ROTATED`, `EXT_SHARE_EXPIRED`.

IP and user-agent hashes use a **daily rotating salt** (`vault_drive_external_share_daily_salts`, pruned after 30 days) — historical hashes become non-reversible after salt rotation for GDPR compliance, while short-term correlation remains available for the owner.

**10.11.5 Key rotation (share rewrap)** The owner can rotate a share at any time without re-uploading the content:

1. `new_shareKey := randomBytes(32)`
2. `new_wrappedFileKey := AES-GCM(fileKey, new_shareKey)`
3. `INSERT new share row; old.linked_to_uuid := new.share_uuid`
4. `old.revoked_at := NOW()`

Recipients using the old link receive “share unavailable”. The owner forwards the new link. All access events from both shares remain linked via `linked_to_uuid` in the audit chain.

**10.11.6 What revocation does — and does not** Revocation **stops future server-side delivery** of the share. It does **not** retroactively revoke share keys, file keys, or chunks already delivered. A recipient who has already downloaded the plaintext — or who captured the link fragment — can continue to use that material locally. For use cases with higher assurance requirements, combine: short `expires_at`, low `max_downloads`, password protection, and the Guardian browser extension for the recipient.

**10.11.7 What the server sees**

| Visible to server                                                   | Not visible                                                        |
|---------------------------------------------------------------------|--------------------------------------------------------------------|
| share_uuid, file_uuid, account_id (owner)                           | Plaintext filename, MIME type, content, message                    |
| wrapped_file_key, wrapped_share_key (ciphertexts)                   | share_key (link-only: lives only in the URL fragment, client-side) |
| unlock_verifier (HMAC output, not reversible)                       | Password, Argon2-derived key                                       |
| protection_mode, expires_at, max_downloads, allow_preview           | Recipient identity (only salted IP/UA hash, pruned after 30 days)  |
| encrypted_filename, encrypted_mime, encrypted_message (ciphertexts) | Their plaintexts                                                   |
| Plaintext sender_display_name (recipient sees it before unlock)     |                                                                    |
| Audit-chain entries for every access event                          |                                                                    |

## 11. Aionda Chat — Post-Quantum E2EE Messaging

Aionda Chat is the integrated real-time messaging surface of Aionda Mail. Unlike email — which uses one-shot ephemeral keys per message — chat conversations are long-lived and require **forward secrecy** and **post-compromise security** for every individual message. To achieve this we designed a dedicated protocol: **AAR (Aionda Async Ratchet)** — a post-quantum variant of Signal’s X3DH + Double Ratchet, where every Diffie–Hellman step is replaced by a hybrid KEM (X25519 + ML-KEM-1024).

### 11.1 Why a Dedicated Protocol?

Email and chat have fundamentally different security profiles:

| Property                        | Email                          | Chat                                                   |
|---------------------------------|--------------------------------|--------------------------------------------------------|
| <b>Cadence</b>                  | Bursty (minutes/hours apart)   | Real-time (seconds apart)                              |
| <b>Session length</b>           | Single message                 | Continuous, days to years                              |
| <b>Forward secrecy unit</b>     | Per message                    | Per message <b>within</b> a long session               |
| <b>Post-compromise recovery</b> | Not required (one-shot key)    | Required — each new message heals from past compromise |
| <b>Asynchrony</b>               | High (recipient often offline) | High (recipient often offline)                         |
| <b>Group dynamics</b>           | Static recipient list          | Dynamic add/remove members                             |

A chat conversation that simply re-used the email pipeline would either (a) burn a fresh ephemeral keypair every keystroke (unacceptably slow), or (b) share a single static conversation key (no forward secrecy). Neither is acceptable. AAR resolves this by maintaining a continuously ratcheting key state per peer — every message advances the state irrevocably.

## 11.2 Building Blocks

AAR uses exactly four primitives:

|            |                      |                                     |
|------------|----------------------|-------------------------------------|
| Hybrid KEM | X25519 + ML-KEM-1024 | (key encapsulation)                 |
| AEAD       | AES-256-GCM          | (message encryption)                |
| KDF        | HKDF-SHA256          | (key derivation)                    |
| Signatures | Ed25519              | (identity-key binding, SPK signing) |

No new cryptographic assumptions are introduced – every primitive is also used elsewhere in the system and has been independently audited.

## 11.3 KeyBundle — The Asynchronous Handshake Material

Every user publishes a **KeyBundle** to the server when chat is first activated. The bundle enables peers to start sending messages even while the user is offline.

KeyBundle (per account, hybrid throughout):

|                         |                                            |
|-------------------------|--------------------------------------------|
| Identity Key (IK)       | – long-lived                               |
| -- IK.x25519_pub        |                                            |
| -- IK.mlkem_pub         |                                            |
| -- IK.ed25519_pub       | – used to sign the SPK                     |
| Signed Pre-Key (SPK)    | – rotated weekly                           |
| -- SPK.x25519_pub       |                                            |
| -- SPK.mlkem_pub        |                                            |
| -- Ed25519 signature    | – signs the concatenation, binds SPK to IK |
| One-Time Pre-Keys (OPK) | – pool of ~100, consumed atomically        |
| -- OPK.x25519_pub       |                                            |
| -- OPK.mlkem_pub        |                                            |

**Server storage:** Only public halves of each key are stored, plus signatures and metadata. Private halves never leave the originating browser. The server enforces an atomic UPDATE ... LIMIT 1 SET consumed\_ts = NOW() when an OPK is fetched, guaranteeing single-use semantics under concurrent requests.

When the OPK pool falls below 20, the client tops up the pool with a fresh batch — this prevents the asynchronous handshake from degrading to a degenerate mode lacking one-time entropy.

## 11.4 PQXDH-Style Handshake (X3DH-PQ)

To start a new conversation with Bob, Alice fetches Bob’s bundle (IK\_B, SPK\_B, one OPK\_B) and Bob’s Ed25519 signature on SPK\_B. Alice verifies the signature, then performs four hybrid KEM encapsulations:

1. Verify Ed25519 signature on SPK\_B – binds SPK to long-term identity
2. Generate ephemeral Alice key (EK\_A):  
EK\_A.x25519, EK\_A.mlkem (one-time, discarded immediately after handshake)
3. Four hybrid KEM encapsulations:  
ssl = HKEM(IK\_A\_priv ⊗ SPK\_B\_pub) – Alice identity → Bob SPK

```

ss2 = HKEM(EK_A_priv ⊗ IK_B_pub) - Alice ephemeral -> Bob identity
ss3 = HKEM(EK_A_priv ⊗ SPK_B_pub) - Alice ephemeral -> Bob SPK
ss4 = HKEM(EK_A_priv ⊗ OPK_B_pub) - Alice ephemeral -> Bob OPK

```

(each `ss_i` is itself the HKDF combination of an X25519 secret AND an ML-KEM secret – see Section 6.4)

4. Root key derivation:

```

SK = HKDF-SHA256(
 IKM = ss1 || ss2 || ss3 || ss4,
 info = "AAR-X3DH-v1" || consumed_OPK_id
)

```

5. Forget every private piece of ephemeral material; SK seeds the ratchet.

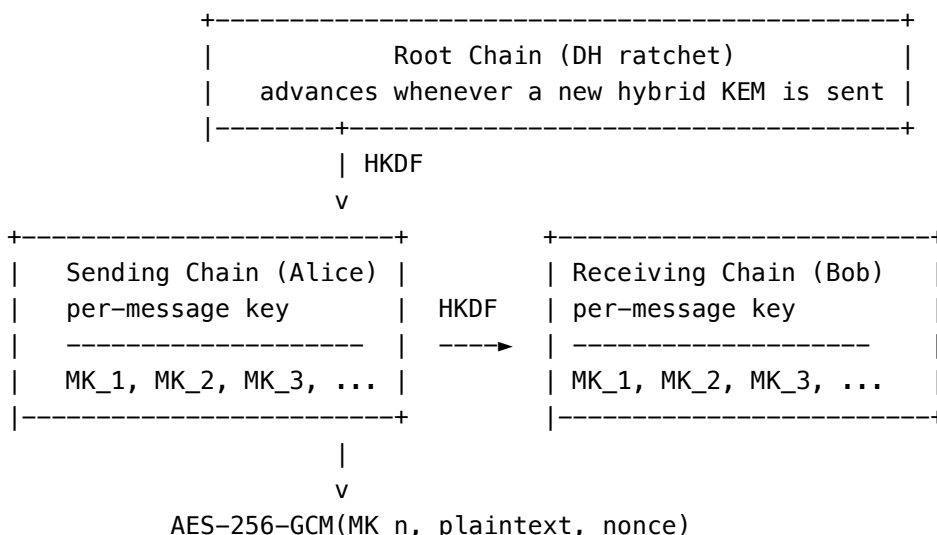
The construction guarantees:

- **Mutual authentication** via the Ed25519-signed SPK\_B and IK\_A binding
- **Forward secrecy** even if IK\_B is later compromised — `ss3` and `ss4` use ephemeral keys on both sides
- **Post-quantum security** because every `ss_i` includes an ML-KEM term — an attacker harvesting today’s traffic for a future quantum computer still cannot reconstruct any of the four secrets
- **Replay resistance** through atomic OPK consumption — Bob’s server-side trigger refuses to release the same OPK twice

The OPK ID consumed by Alice is bound into the `info` parameter of HKDF — Bob’s side can therefore reproduce the same root key only with the same OPK, and only once.

### 11.5 Double Ratchet — Per-Message Keys

After the handshake, both sides advance a **Double Ratchet** state for every message:



Every chat event carries the sender’s current **ratchet public key** (hybrid pair, ~1.6 KB). When the recipient receives a message whose ratchet key differs from the last one seen, both parties perform a fresh hybrid KEM, derive a new root key, and reset their chains. This is the **DH ratchet**

in Signal terminology — except every step is a full hybrid KEM, not a single Diffie–Hellman.

### Per-message guarantees:

- **Forward secrecy:** Deletion of `MK_n` makes message `n` undecryptable, even given the full long-term state at the time of capture
- **Post-compromise security:** Once a fresh DH-ratchet step has occurred after a compromise, all subsequent messages are protected from the attacker
- **Out-of-order delivery:** Message keys are cached for late-arriving messages (capped at 1000 keys per session to bound memory)

## 11.6 Message Container

Every chat event uploaded to the server is a self-contained ciphertext:

MessageContainer (base64-encoded, stored in `chat_events.payload`):

```
+-----+
| ratchet_pub : Hybrid KEM public key (current step) |
| prev_chain_len : Number of messages in previous chain |
| msg_number : Counter within current chain |
| ciphertext : AES-256-GCM(MK_n, plaintext || header) |
| tag : AES-256-GCM authentication tag |
+-----+
```

The header inside the AEAD plaintext contains: sender login string, `content_type` (`text/quote/file-reference`), `timestamp`, optional `reply_to_event_uuid`.

The container is opaque to the server — including the message type, the sender’s display name format, and any quoted reply context. Aionda’s database stores the container as-is in `chat_events.payload`.

## 11.7 Group Conversations — Per-Recipient Sender Keys

For room conversations (3+ participants), AAR uses a **per-recipient fan-out** model. The sender derives one sender-key chain per peer pair and encrypts a message once per recipient. Each recipient therefore receives a personally-addressed ciphertext, decryptable only with the keys derived from their own AAR session with the sender.

**Trade-off:** Fan-out cost is  $O(N)$  in recipients — acceptable for typical team sizes ( $\leq 25$  participants). For larger groups a future MLS-based mode is planned (see Roadmap). The current model is preferable to a single shared group key because: (a) it preserves forward secrecy on a per-pair basis, (b) it does not require key-rotation on member removal beyond the bilateral level, and (c) it inherits the post-quantum guarantees of pairwise AAR without modification.

## 11.8 Initial Conversation Container

When Alice starts a new room with Bob, Carol, and Dan, the first message to each participant is wrapped as an **initial-handshake container** that carries (a) the X3DH handshake material consumed against that participant’s bundle and (b) the first message itself. The receiver dispatches on the discriminator field:

```
{
 "type": "initial",
 "initialMessage": { /* handshake metadata, OPK id consumed */,
 "encryptedMessage": "<base64 MessageContainer>"
}
```

Subsequent messages in the same room reuse the established AAR session and only contain the encryptedMessage portion.

### 11.9 Transport — Encrypted API + Mercure SSE + WebSocket

The chat plane uses three orthogonal channels:

| Channel                   | Direction        | Purpose                                                              | Plaintext seen by server                                                               |
|---------------------------|------------------|----------------------------------------------------------------------|----------------------------------------------------------------------------------------|
| <b>/e Encrypted API</b>   | Client -> Server | Send message, fetch history, key-bundle ops (11 endpoints)           | None — all 11 chat endpoints use the same hybrid-KEM transport documented in Section 8 |
| <b>Mercure SSE</b>        | Server -> Client | Push delivery of new events (chat.event_received, chat.read_receipt) | None — the server pushes the same opaque MessageContainer it stored                    |
| <b>WebSocket /chat/ws</b> | Bi-directional   | Reconnect backfill (sync_request), heartbeat, presence               | Presence boolean (online/offline) and login-string only — never message content        |

Presence is held in an in-memory map on `aionda_chat_realtime`; a single intranet-only HTTP endpoint exposes a `login -> online` lookup for the team picker so users can see whether a peer is reachable. No history, no read state, no content ever traverses the presence channel.

### 11.10 Read Receipts

Read receipts are an opt-in per-account setting (`chat_participants.send_read_receipts`). When enabled, marking a message as read publishes a `chat.read_receipt` event containing:

```
{ conversation_uuid, last_read_event_uuid, account_login }
```

Note that `last_read_event_uuid` is **not** the message content — it is a server-allocated identifier already known to the server. No additional information leaks beyond “Alice has now seen messages up to event X.” Recipients who disable read receipts (`send_read_receipts = false`) never emit such events, and the server enforces the toggle.

### 11.11 Audit Chain Integration

For Enterprise accounts, every chat operation is appended to the existing tamper-evident audit log (`enterprise_audit_log`, SHA3-256 hash chain — see §18). The reserved action types are:

```
CHAT_CONVERSATION_CREATED, CHAT_PARTICIPANT_ADDED, CHAT_PARTICIPANT_REMOVED, CHAT_MESSAGE_SENT,
```

CHAT\_MESSAGE\_READ, CHAT\_KEYBUNDLE\_PUBLISHED, CHAT\_KEYBUNDLE\_ROTATED, CHAT\_OPK\_TOPPED\_UP, CHAT\_CONVERSATION\_LEFT.

The audit record contains only the conversation UUID, the actor's login, and a timestamp — **never** the encrypted payload, ratchet keys, or message contents.

### 11.12 What the Server Sees

| Visible to server                                              | Not visible                                                                                                          |
|----------------------------------------------------------------|----------------------------------------------------------------------------------------------------------------------|
| conversation_uuid, list of participants (by mail_account.name) | Plaintext message content                                                                                            |
| Encrypted MessageContainer payloads (opaque ciphertexts)       | Ratchet keys, chain keys, message keys                                                                               |
| Public halves of IK, SPK, OPK; Ed25519 signatures on SPK       | Private halves of any keypair                                                                                        |
| OPK consumption counter and timestamp                          | Which OPK was consumed for which conversation (correlation prevented by HKDF-info binding only existing client-side) |
| Conversation cadence (timestamps of chat_events rows)          | Quoted reply chains, file attachments inside the container                                                           |
| Presence boolean (online/offline)                              | Typing indicators (peer-to-peer only, never reach the server)                                                        |
| Audit-chain entries for Enterprise accounts                    | Plaintext content of audit entries                                                                                   |

### 11.13 Cryptographic Audits

The AAR protocol was implemented from scratch in TypeScript and underwent three independent AI-assisted audits documented in typescript/manager/chat/crypto/AUDIT.md, AUDIT\_SECOND\_OPINION.md, and AUDIT\_THIRD\_OPINION.md. Findings from each round were incorporated before public deployment. The primary surfaces audited were: handshake symmetry, OPK reservation under concurrency, deep-cloning of immutable ratchet state, and the AEAD framing of the MessageContainer.

### 11.14 Implementation Boundaries

The chat code base lives in typescript/manager/chat/ (~9,000 lines of TypeScript) and includes/classes/App (PHP service layer). Key audit artifacts:

```
crypto/aar-types.ts - type-only declarations, no logic
crypto/aar-keybundle.ts - KeyBundle creation, rotation, serialization
crypto/aar-x3dh.ts - Handshake (initiator + responder paths)
crypto/aar-double-ratchet.ts - Root chain, sending chain, receiving chain
```

The clean separation between protocol logic (crypto/) and transport/UI (chat-store.ts, chat-panel.ts, ...) ensures that the cryptographic core can be re-audited without UI scope creep.

## 12. Aionda Video Calls — Post-Quantum E2EE Conferencing

Aionda Mail includes end-to-end encrypted group video calls (cipher suite AIONDA\_PQ\_CALL\_V1). The media relay never sees plaintext audio or video — it forwards only ciphertext frames. Group keys are agreed with a **post-quantum hybrid** handshake.

## 12.1 Security Goals & Trust Model

### In scope — the protocol provides:

- **Confidentiality** of audio and video media, end to end.
- **Authenticity** of the group — every participant is cryptographically a member of the same MLS group, verified out-of-band through the Short Authentication String (Section 12.5).
- **Forward secrecy** — past media stays secret even if current keys are later compromised.
- **Post-compromise security** — the group recovers confidentiality after a compromised member is removed and the group rekeys.
- **Post-quantum confidentiality** — group establishment uses the X-Wing hybrid KEM.
- **Server independence** — a malicious or compromised SFU, signaling server or storage layer cannot read or inject media.

### Out of scope (explicitly *not* defended against):

- Traffic analysis — packet timing, packet sizes, participant count and call duration are observable (Section 12.9).
- Endpoint compromise (see trust assumption below).
- Operating-system-level malware or physical device seizure.

**Endpoint trust assumption.** The model assumes an uncompromised browser and operating system. Frames are sealed inside the browser media pipeline (`RTCRtpScriptTransform`), but a compromised endpoint can read plaintext media *before* encryption or *after* decryption. The web-delivery trust model and its mitigations (Guardian response signing) are covered in Section 17 and Section 22.

## 12.2 Why a Dedicated Protocol?

Real-time conferencing has two requirements that the email and chat layers do not:

1. **Per-frame encryption at line rate.** Audio/video frames must be encrypted individually so a Selective Forwarding Unit (SFU) can route, drop and re-order them without ever decrypting. This is the role of **SFrame** (RFC 9605).
2. **Dynamic group key agreement.** Participants join and leave mid-call. Each membership change must rekey the group with forward secrecy and post-compromise security. This is the role of **MLS** (RFC 9420).

The server-side SFU (`mediasoup`) is treated as **untrusted infrastructure**: it is a zero-knowledge packet router, exactly like the storage layer for email.

## 12.3 Cipher Suite — `AIONDA_PQ_CALL_V1 (0xFF01)`

The MLS profile substitutes the classical primitives of a standard MLS suite with post-quantum hybrids.

| Role       | Algorithm                           | Notes                                    |
|------------|-------------------------------------|------------------------------------------|
| HPKE KEM   | <b>X-Wing</b> (ML-KEM-768 + X25519) | Hybrid;<br>draft-connelly-cfrg-xwing-kem |
| Signatures | <b>ML-DSA-65</b>                    | NIST FIPS 204 (Dilithium),<br>Level 3    |
| KDF        | HKDF-SHA-256                        | RFC 5869                                 |

| Role                 | Algorithm          | Notes                                                                               |
|----------------------|--------------------|-------------------------------------------------------------------------------------|
| AEAD (HPKE + SFrame) | <b>AES-256-GCM</b> | AES-256 provides a comfortable margin against known quantum search (Grover) attacks |

**No downgrade.** There is no classical-only fallback. If the post-quantum handshake cannot be completed, the call fails closed — a downgrade is a hard error, never a silent weakening.

**X-Wing maturity.** X-Wing is currently an IETF/CFRG draft (`draft-connelly-cfrg-xwing-kem`). The current profile adopts it as the preferred hybrid KEM; future protocol revisions may migrate to a finalized CFRG/NIST-standardized hybrid KEM while preserving the MLS and SFrame layers unchanged.

**Cryptographic agility.** The cipher-suite identifier (`0xFF01`) fully determines the KEM, signature algorithm, KDF and AEAD. A future `AIONDA_PQ_CALL_V2` can replace any primitive without changing the higher protocol layers; suites are selected by identifier and never silently negotiated downward.

#### 12.4 Group Key Agreement (MLS, RFC 9420)

Every participant is represented in the MLS ratchet tree by a **KeyPackage** (X-Wing public KEM key + ML-DSA-65 credential). The group evolves through *epochs*:

- **Welcome / Commit** — a joining member is added with an X-Wing-encapsulated path secret; the ratchet tree advances to a new epoch.
- **Rekeying on every membership change.** Any change — a join, a leave, a device replacement, or an identity-key rotation — advances the group to a fresh epoch. This is what delivers forward secrecy (a new joiner cannot read past media) and post-compromise security (a removed member cannot read future media).
- **Remove-Commit.** When a participant leaves (or a stale/zombie peer is detected), the owner issues a Remove-Commit that rekeys the group, so a departed member cannot decrypt subsequent frames even if it retained old key material.

The hybrid construction is designed so that compromising only **one** component — ML-KEM-768 or X25519 — does not by itself reveal the group secret.

#### 12.5 Identity & Authentication

- **Credentials.** Each member carries an MLS **Credential** that binds a participant identity to its ML-DSA-65 signature key. Every MLS handshake message (Welcome / Commit / Update) is signed with that key, so the SFU or signaling server cannot forge group operations.
- **Short Authentication String (SAS).** The room link carries a high-entropy **pre-shared secret** in the URL fragment, which is never sent to the server in the clear. The MLS group exporter and this secret are folded together into a Short Authentication String, whose fingerprint the client **pins**. A mismatch — for example a signaling server or SFU attempting to splice in a rogue participant — triggers a **hard abort with no soft fallback**. This binds the cryptographic group to the out-of-band room secret and defeats server-side man-in-the-middle.
- **Authentication factor.** For a guest room, *possession of the room link is equivalent to*

*possession of the room credential* — the room secret is the sole authentication factor. Authenticated Aionda participants additionally redeem a signed, account-scoped signaling token (Section 12.8).

## 12.6 Media Encryption (SFrame, RFC 9605)

Each media track is encrypted frame-by-frame **before** it reaches the SFU. The key lifecycle:

```

MLS group secret
 | MLS Exporter (RFC 9420 §8)
 v
SFrame base key
 | HKDF(base_key, KID = sender || epoch)
 v
per-sender key
 | AES-256-GCM(nonce = frame counter, aad = sframe_header)
 v
per-frame ciphertext || tag -> wire = sframe_header || ciphertext || tag

```

- The SFrame key is derived from the **MLS Exporter**, so it exists only on authenticated group members — never on the server.
- Encryption runs in a dedicated `RTCRtpScriptTransformWorker` (the standardised Encoded-Transform API, available in Chromium and Safari/WebKit), so frames are sealed on the browser’s media pipeline.
- The codec is VP8; the SFU performs keyframe routing on metadata only — it never inspects the encrypted payload.

## 12.7 Replay Protection

Each encrypted frame carries a monotonic SFrame **counter** inside its authenticated header, and is bound to the current MLS epoch. Receivers track the highest counter seen per sender (a persisted receive-side counter store). A frame with a duplicate or regressed counter, or one bound to a stale epoch, is rejected — a captured ciphertext cannot be replayed into the session.

## 12.8 Signaling & Access Model

- **Signaling** runs over an authenticated WebSocket (`aionda_signaling`). Each session is gated by a short-lived **JWT** minted by the application server after the client redeems the room secret via the encrypted `/e` API (see Section 8). The token is symmetric (HS256) and validated only between the application server and the single signaling service; it authorizes session setup and never carries media or group secrets. The signaling server relays MLS Welcome/Commit messages but cannot read them.
- **Room link = bearer credential.** Anyone holding the full link (room UUID + fragment secret) may join — the model is intentionally “whoever has the link”. Database IDs are never exposed; rooms are addressed only by unguessable UUIDs.
- **Owner control.** The room owner (verified server-side by account ownership) may mark a room *ended* to block future joins, and may remove a present participant via an MLS Remove-Commit (Section 12.4).

## 12.9 What the SFU / Server Sees — and Metadata Limitations

| The infrastructure CAN see                      | The infrastructure CANNOT see                                            |
|-------------------------------------------------|--------------------------------------------------------------------------|
| Encrypted frame sizes, timing, routing metadata | Audio or video content (frames are AES-256-GCM sealed)                   |
| MLS handshake ciphertext (Welcome/Commit)       | Group secrets or SFrame keys (derived client-side from the MLS exporter) |
| Room UUID, participant count, JWT subject       | The room link secret (lives in the URL fragment)                         |
| ICE candidates / transport addresses            | Any plaintext that would let it join or decrypt the session              |

**Metadata the system intentionally does not hide:** IP addresses exposed to ICE / STUN / TURN infrastructure, packet timing, packet sizes, participant count, and call duration. Traffic-analysis resistance is not a design goal (Section 12.1).

### 12.10 Security Properties (Summary)

| Property                 | Guarantee                                                                                             |
|--------------------------|-------------------------------------------------------------------------------------------------------|
| Confidentiality          | Only current MLS group members possess the SFrame keys.                                               |
| Authentication           | All MLS messages are signed with ML-DSA-65; the group is bound to the room secret via the pinned SAS. |
| Forward secrecy          | Compromise of current secrets does not reveal past media.                                             |
| Post-compromise security | Rekeying restores confidentiality after a member is removed.                                          |
| Server independence      | Compromise of the signaling server, SFU or storage does not expose media.                             |
| Post-quantum             | Group establishment relies on the X-Wing hybrid KEM (ML-KEM-768 + X25519).                            |

These properties describe the protocol *design*. An independent third-party cryptographic audit of the call subsystem is planned; until it is published, the guarantees above are a self-assessment.

## 13. Side-Channel Protections

### 13.1 Bucket Padding

**Problem:** Encrypted email sizes can leak information. An attacker observing ciphertext lengths could infer content (e.g., a 50-byte email is likely “OK, thanks” while a 500KB email contains attachments).

**Solution:** All data is padded to fixed-size “buckets” before encryption:

Bucket sizes: 256B, 512B, 1KB, 2KB, 4KB, 8KB, 16KB, 32KB,  
64KB, 128KB, 256KB, 512KB, 1MB, 2MB, 4MB, 8MB, 16MB

Format: [0xDEAD magic][4-byte length][actual data][random padding to bucket boundary]

**Example:** A 523-byte email is padded to 1,024 bytes. An observer sees only “1KB email” — not the actual 523-byte size.

### 13.2 Compression-Before-Encryption

Data is compressed with gzip (level 6) **before** encryption. This is the only correct order:

- Compression after encryption would fail (encrypted data has maximum entropy)
- The bucket padding after compression prevents CRIME/BREACH-style attacks that exploit compression ratios

### 13.3 Threading Privacy

Email threads use SHA-256 hashes of Message-ID headers instead of plaintext identifiers. The server can group related emails by hash equality without knowing the actual message identifiers.

---

## 14. Key Management & Lifecycle

### 14.1 Key Hierarchy

Vault Master Key (32 bytes, generated once per account)

```

|
|--- Vault Keypair (Hybrid KEM)
| |-- X25519 public key (32 bytes) – stored plaintext on server
| |-- X25519 private key (32 bytes) – encrypted with master key
| |-- ML-KEM-1024 public key (1568 bytes) – stored plaintext on server
| |-- ML-KEM-1024 private key (3168 bytes) – encrypted with master key
|
|--- Per-Email Ephemeral Keys (32 bytes each)
| |-- Wrapped with recipient's Hybrid KEM public keys
|
|--- Per-Attachment Ephemeral Keys (32 bytes each)
| |-- Wrapped independently per attachment
|
|--- Folder Keys (derived via HKDF per folder)
| |-- Shared via Hybrid KEM encapsulation per recipient
|
|--- Signature Encryption Key (derived from master key)
| |-- Encrypts email signature templates

```

### 14.2 Key Storage

| Key                | Storage Location                                     | Protection                  |
|--------------------|------------------------------------------------------|-----------------------------|
| Master Key         | Nowhere (reconstructed on-demand from Shamir shares) | Shamir 2-of-3               |
| Vault private keys | Server (encrypted)                                   | AES-256-GCM with master key |

| Key                     | Storage Location            | Protection                                      |
|-------------------------|-----------------------------|-------------------------------------------------|
| Vault public keys       | Server (plaintext)          | Not sensitive — public by definition            |
| Email ephemeral keys    | Server (wrapped)            | Hybrid KEM encapsulation                        |
| OPAQUE records          | Server (encrypted at rest)  | AES-256-GCM with server key                     |
| Encrypted Shamir shares | Server                      | XOR with password/passkey/recovery derived keys |
| API transport keys      | Server (pre-generated pool) | One-time use, 24h TTL                           |

### 14.3 Key Fingerprints

Each vault keypair has a SHA-256 fingerprint stored on the server. This allows:

- Audit trail of key rotations
- Detection of unauthorized key changes
- Client-side verification of key integrity

## 15. Recovery Mechanism

### 15.1 Recovery Key (BIP39 Mnemonic)

During vault setup, the user is presented with a **24-word recovery phrase** generated from 256 bits of entropy, encoded using the BIP39 standard:

Example: apple river mountain sunset golden bridge falcon ocean  
 crystal thunder meadow silver dolphin forest marble castle  
 velvet compass harbor window ancient pepper rocket shield

### 15.2 Recovery Key Derivation

1. Generate: 256 bits random entropy
2. Encode: BIP39 mnemonic (24 words, 11 bits per word)
3. Derive: `verificationKey = HKDF-SHA3-256(entropy, salt = accountId, info = "trashmail-recovery-verify")`
4. Hash: `verificationHash = SHA3-256(verificationKey)`
5. Store: Server stores **ONLY** verificationHash (32 bytes)

### 15.3 What the Server Stores

The server stores **only the SHA3-256 hash** of a derived verification key. It does not store:

- The recovery words

- The entropy
- The verification key itself

#### 15.3.4 Recovery Flow

1. User enters 24-word recovery phrase
2. Client derives entropy -> HKDF-SHA3-256 -> SHA3-256 -> verification hash
3. Client sends verification hash to server (never the plaintext key)
4. Server compares with stored hash
5. If match: All 2FA methods are disabled, user sets up fresh authentication
6. Recovery key is revoked after single use

#### 15.3.5 Rate Limiting

- Maximum 3 verification attempts per hour
- 60-minute lockout after exceeding limit
- One-time use: recovery key is permanently revoked after successful use

#### 15.3.6 No Password Recovery

**There is no password reset via email.** If a user loses their password AND all other authentication factors (passkey + recovery key), their data is permanently inaccessible. This is the fundamental proof that zero-knowledge works — if the service could recover user data, it could also read it.

---

## 16. Passwordless Authentication (Passkeys)

### 16.1 WebAuthn PRF Extension

Aionda Mail supports FIDO2 passkeys (hardware security keys, biometric authenticators) for passwordless login and vault unlock.

The **WebAuthn PRF (Pseudo-Random Function) extension** provides a deterministic 32-byte output bound to the specific passkey and credential. This output is used to protect Shamir Share 2.

### 16.2 How It Works

1. Registration:
  - User creates passkey via `navigator.credentials.create()`
  - PRF extension generates hardware-bound output
  - Output XOR'd with Shamir Share 2 -> encrypted share stored on server
2. Authentication:
  - User authenticates with passkey (biometric/PIN)
  - PRF extension reproduces same 32-byte output
  - Output XOR'd with encrypted share -> Shamir Share 2 recovered
  - Combined with Share 1 (password) -> Master Key reconstructed
3. Vault Unlock (passwordless):
  - If both passkey (Share 2) and password (Share 1) available -> immediate unlock

- Password verified via OPAQUE (separate from passkey auth)

### 16.3 Multiple Passkeys

Users can register multiple passkeys (e.g., MacBook Touch ID, iPhone Face ID, YubiKey). Each passkey independently protects its own copy of Share 2. Any single passkey combined with the password is sufficient to unlock the vault.

## 17. Guardian: MITM Protection & Response Signing

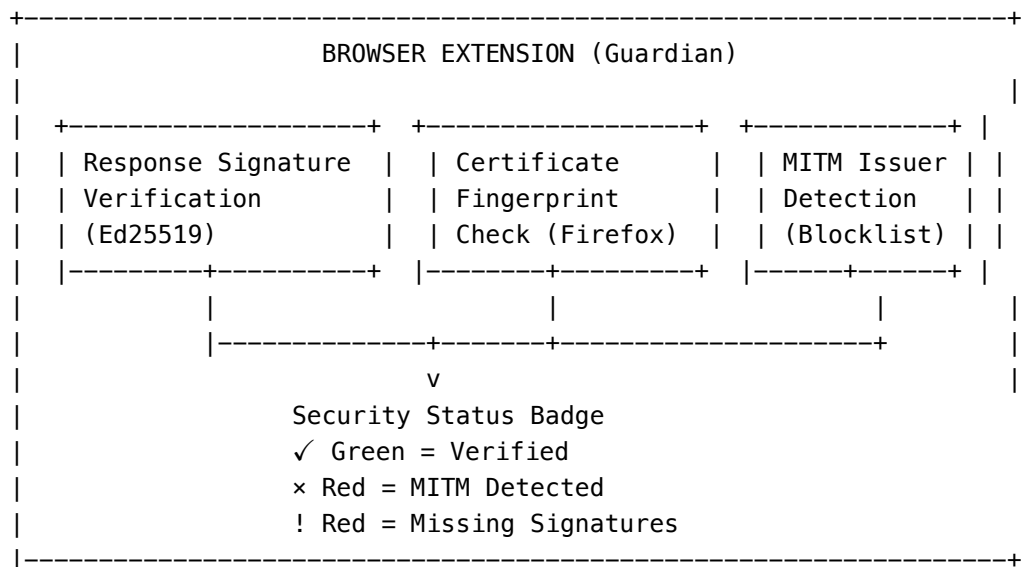
### 17.1 The Problem with Web Applications

Every web application has an inherent trust problem: the browser downloads JavaScript from the server on every visit. A man-in-the-middle (MITM) attacker — whether a compromised CDN, corporate proxy, or rogue ISP — could theoretically inject modified code that exfiltrates encryption keys.

Aionda Mail addresses this with the **Guardian module**, a browser extension component (available for Chrome and Firefox) that independently verifies server integrity.

### 17.2 Architecture Overview

The Guardian module operates inside the browser extension’s Service Worker — completely independent from the web application’s JavaScript. It performs three types of verification:



### 17.3 Response Signature Verification (Ed25519)

Every API response from Aionda Mail’s server is cryptographically signed using **Ed25519** (Edwards-curve Digital Signature Algorithm).

#### Signing process (server-side):

1. Server generates API response body (JSON)
2. Construct signing input: responseBody + "|" + unixTimestamp
3. Sign with Ed25519 private key -> 64-byte signature
4. Attach HTTP headers:

```
X-Aionda-Signature: <base64(signature)>
X-Aionda-Timestamp: <unix_timestamp>
X-Aionda-Key-Id: <key_identifier>
```

### Verification process (browser extension):

1. Extract signature, timestamp, and key ID from HTTP headers
2. Look up Ed25519 public key by key ID (bundled in extension)
3. Verify key has not expired (valid\_from / valid\_until)
4. Check timestamp freshness:  $|\text{now} - \text{timestamp}| \leq 300$  seconds
5. Reconstruct signing input: `responseBody + "|" + timestamp`
6. `crypto.subtle.verify("Ed25519", publicKey, signature, data)`
7. If invalid -> MITM alert, red badge

### Key properties:

- **Replay protection:** 5-minute timestamp window prevents replaying old responses
- **Tamper detection:** Any modification to the response body invalidates the signature
- **Key isolation:** Public keys are bundled inside the extension (not downloaded from the server)
- **Environment separation:** Dev keys (dev-2026-01) cannot be used on production URLs and vice versa

## 17.4 Ed25519 Public Key Management

Public keys are shipped with the browser extension in `public_key.json`:

```
{
 "keys": {
 "prod-2026-01": {
 "algorithm": "Ed25519",
 "public_key": "<base64 SPKI DER>",
 "valid_from": "2026-01-13T00:00:00Z",
 "valid_until": "2027-01-13T00:00:00Z"
 }
 }
}
```

- **Key format:** SPKI DER (Subject Public Key Info, Distinguished Encoding Rules)
- **Key size:** 32 bytes (256-bit Ed25519 public key)
- **Signature size:** 64 bytes (fixed)
- **Rotation:** New keys are added before old keys expire; extension updates deliver new keys
- **No server trust:** Keys are embedded in the extension binary, not fetched from the server

## 17.5 TLS Certificate Verification (Firefox)

On Firefox, the Guardian module performs additional TLS certificate verification using the `browser.webRequest.getSecurityInfo()` API (not available in Chrome due to Manifest V3 limitations).

### Verification flow:

1. Browser extension intercepts HTTPS response
2. Extract TLS certificate chain from browser's security info:

- Leaf certificate fingerprint (SHA-256)
  - Issuer Distinguished Name (O=, CN=)
  - Subject (CN=)
3. Check against known MITM issuers (hardcoded blacklist):  
ZScaler, Netskope, Fortinet, Palo Alto, Blue Coat,  
Check Point, Barracuda, Sophos, WatchGuard, Cisco Umbrella  
-> If match: MITM detected, show warning
  4. Check against trusted issuers:  
Google Trust Services, Cloudflare, Let's Encrypt,  
DigiCert, Sectigo  
-> If match AND subject matches expected domain: OK
  5. If unknown issuer: Fetch server's own certificate fingerprint
    - Server connects to itself via external routing (prevents spoofing)
    - Response is Ed25519 signed (prevents MITM from lying about cert)
    - Compare issuer organization with browser's certificate issuer
    - > If mismatch: MITM suspected, show warning

**Why issuer-based validation instead of pinning?** CloudFlare (used as CDN) rotates leaf certificates across edge servers. Traditional certificate pinning (matching exact fingerprints) would cause false positives. Issuer-based validation is more robust: the issuing CA is stable even when leaf certificates change.

### 17.6 Self-Certificate Fetching (Anti-Spoofing)

The server's certificate endpoint uses a clever anti-spoofing technique:

Server connects to cert.trashmail.com (or cert-subdomain.domain)  
with SNI = mail.aionda.com

- > Forces external routing through CloudFlare
- > Receives the actual certificate that users see
- > Prevents localhost spoofing
- > Response signed with Ed25519 to prevent tampering

The server essentially asks “what certificate does the outside world see for my domain?” — and signs the answer so the extension can trust it.

### 17.7 Security Status Indicators

The extension displays a badge on the browser toolbar:

| Badge    | Color  | Meaning                                         |
|----------|--------|-------------------------------------------------|
| ✓        | Green  | All responses verified — signatures valid       |
| !        | Orange | Using deprecated signing key (rotation pending) |
| ×        | Red    | MITM detected — signature verification failed   |
| !        | Red    | Missing signatures — responses not signed       |
| [Shield] | Blue   | Protected mode — no verification performed yet  |

### 17.8 Threat Coverage

| Attack                                   | Detection Method                      | Browser          |
|------------------------------------------|---------------------------------------|------------------|
| Corporate MITM proxy (ZScaler, Fortinet) | Certificate issuer blocklist          | Firefox          |
| Modified API responses                   | Ed25519 signature verification        | Chrome + Firefox |
| Replay attacks                           | 5-minute timestamp window             | Chrome + Firefox |
| CDN compromise (CloudFlare)              | Response signature mismatch           | Chrome + Firefox |
| Certificate substitution                 | Issuer comparison + server self-check | Firefox          |
| Dev/prod key confusion                   | Environment-bound key IDs             | Chrome + Firefox |

### 17.9 Limitations

- **Chrome Manifest V3:** Cannot inspect TLS certificates — only response signature verification is available
- **Extension required:** Users without the extension do not benefit from Guardian protections
- **Ed25519 is not post-quantum:** Signature verification uses classical cryptography. A sufficiently powerful quantum computer could theoretically forge Ed25519 signatures.

## 18. Enterprise Email Archive (Blockchain)

### 18.1 Overview

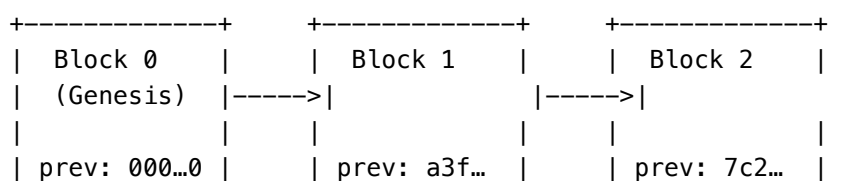
Aionda Mail’s Enterprise plan includes a **GoBD-compliant email archive** secured by a cryptographic hash chain (blockchain). Every archived email becomes an immutable block in a per-company chain. Any tampering — modification, deletion, or insertion of blocks — is cryptographically detectable.

The archive combines two independent security layers:

1. **Hash chain (SHA3-256):** Guarantees integrity and immutability — proves no email was altered or removed after archival
2. **Hybrid KEM encryption (CAK):** Guarantees confidentiality — the server cannot read archived email content

### 18.2 Hash Chain Architecture

Each archived email becomes a block in a sequential, tamper-evident chain:



```
| hash: a3f... | | hash: 7c2... | | hash: e91... |
| email: b17... | | email: 4d9... | | email: f83... |
|-----+ |-----+ |-----+
```

**Block hash calculation:**

```
block_hash = SHA3-256(
 prev_block_hash || "|" ||
 timestamp || "|" ||
 email_hash || "|" ||
 direction || "|" ||
 sender_domain || "|" ||
 recipient_domain
)
```

**Properties:**

- **Hash algorithm:** SHA3-256 (NIST FIPS 202)
- **Genesis block:** prev\_block\_hash = 64 zeros, block\_number = 0
- **Sequential numbering:** Enforced by database UNIQUE KEY (company\_uuid, block\_number)
- **One chain per company:** Complete isolation between enterprises
- **Email hash:** SHA3-256(sender || recipient || timestamp || size) — integrity proof of the original email data

**18.3 Tamper Detection**

The chain verification algorithm detects any form of tampering:

For each block (ordered by block\_number ASC):

1. Verify link:     block.prev\_block\_hash == expected\_prev\_hash
2. Recalculate:    expected = SHA3-256(prev\_hash | timestamp | email\_hash | ...)
3. Verify content: block.block\_hash == expected
4. Advance:        expected\_prev\_hash = block.block\_hash

If ANY check fails -> chain is broken at block N

| Tampering Attempt                           | Detection                                                           |
|---------------------------------------------|---------------------------------------------------------------------|
| Modify email content                        | email_hash changes -><br>block_hash recalculation fails             |
| Modify metadata (sender, domain, timestamp) | Included in hash input -><br>block_hash mismatch                    |
| Delete a block                              | Next block's prev_block_hash<br>becomes orphaned                    |
| Insert a block                              | Breaks sequential block_number<br>+ prev_block_hash chain           |
| Reorder blocks                              | UNIQUE KEY constraint +<br>sequential verification prevents<br>this |
| Replace entire chain                        | Genesis block hash would differ<br>from any external backup         |



```
"ha": true,
"ac": 3,
"en": "John Doe"
}
```

**Encrypted email content** (separate Hybrid KEM key wrapping):

```
{
 "subject": "Meeting notes",
 "body": "<html>...</html>",
 "from": "sender@domain.com",
 "to": "recipient@company.de"
}
```

**Zero-knowledge enforcement:** After encryption, the plaintext metadata fields in the database (sender\_address, recipient\_address, domains) are replaced with their SHA3-256 hashes. The server stores only hashes — the original values exist only inside the encrypted blobs.

## 18.6 Audit Trail

Every action on the archive is logged in an **independent audit chain** (also hash-chained with SHA3-256):

| Action                                                   | When Logged                   |
|----------------------------------------------------------|-------------------------------|
| EMAIL_RECEIVED / EMAIL_SENT / DRAFT_ARCHIVED             | Email archived                |
| VIEW_EMAIL / VIEW_ATTACHMENT                             | Employee reads archived email |
| SEARCH_ARCHIVE                                           | Search performed              |
| EXPORT_EMAIL / EXPORT_REPORT                             | Data exported                 |
| VERIFY_CHAIN / CHAIN_VERIFIED_OK / CHAIN_VERIFIED_BROKEN | Integrity check               |
| LEGAL_HOLD_SET / LEGAL_HOLD_RELEASED                     | Legal hold toggled            |
| ARCHIVE_DECRYPT                                          | CAK used to decrypt content   |
| ADMIN_ACCESS                                             | Administrative action         |

Each audit entry records: actor (UUID + role), IP address, session ID, target email hash, and whether the chain was valid at the time of access.

## 18.7 Legal Hold & Retention

- **Retention period:** Configurable per company (default: 10 years), calculated per email as archived\_at + retention\_years
- **Legal hold:** Individual emails can be placed under legal hold, preventing deletion until released. Includes reason, actor, and timestamp.
- **GoBD compliance:** The combination of immutable hash chain, complete audit trail, configurable retention, and legal hold satisfies the requirements of the German GoBD (Grundsätze zur ordnungsmäßigen Führung und Aufbewahrung von Büchern, Aufzeichnungen und Unterlagen in elektronischer Form sowie zum Datenzugriff).

## 18.8 Forensic Export

Authorized users (Owner, Admin) can export the complete chain state for independent verification:

- Full chain data with all block hashes
- Verification result (valid/broken, broken block number if applicable)
- Expected vs. actual hash values for forensic analysis
- Last 50 audit log entries
- JSON format for external re-verification with any SHA3-256 implementation

## 19. What the Server Sees — and What It Does Not

This section explicitly documents the zero-knowledge boundary.

### 19.1 The Server CAN See

| Data                    | Why Visible                       | Mitigation                                   |
|-------------------------|-----------------------------------|----------------------------------------------|
| IP address              | TCP/IP requirement                | Use VPN/Tor if desired                       |
| Timestamps              | Email reception time              | Inherent to email protocol                   |
| Encrypted email blobs   | Stored for retrieval              | AES-256-GCM encrypted, key unknown to server |
| Padded ciphertext sizes | Storage requirement               | Bucket padding hides actual sizes            |
| Recipient DEA address   | Routing requirement               | DEA is disposable, not the real address      |
| Account existence       | Authentication flow               | User enumeration protection deployed         |
| Public keys             | Required for encryption by server | Public by definition, not sensitive          |
| Encrypted Shamir shares | Storage for user                  | XOR'd with keys server doesn't know          |
| OPAQUE records          | Authentication protocol           | Not password hashes, encrypted at rest       |

### 19.2 The Server CANNOT See

| Data                                   | Why Invisible                                        |
|----------------------------------------|------------------------------------------------------|
| Email content (subject, body, headers) | Encrypted with ephemeral keys wrapped via Hybrid KEM |
| User password                          | OPAQUE — password never transmitted                  |
| Master key                             | Reconstructed only in browser from Shamir shares     |
| Vault private keys                     | Encrypted with master key before storage             |
| Email ephemeral keys                   | Wrapped with Hybrid KEM, server lacks private keys   |
| Recovery key / mnemonic                | Only SHA3-256 hash of derived key stored             |
| Passkey PRF outputs                    | Hardware-bound, never leave authenticator            |
| Folder names                           | Encrypted with folder-specific keys                  |
| Email signatures                       | Encrypted with master key                            |

| Data                 | Why Invisible                    |
|----------------------|----------------------------------|
| API request content  | Encrypted via /e transport layer |
| API response content | Encrypted before transmission    |

### 19.3 Cryptographic Guarantee

Even with full access to:

- The complete database
- All network traffic
- The server's source code and configuration
- All OPAQUE records and server keys

...an attacker **cannot** decrypt a single email without the user's password (or passkey + recovery key). This is not a policy — it is a mathematical impossibility enforced by the cryptographic design.

## 20. Algorithm Reference

### 20.1 Complete Algorithm Table

| Component                 | Algorithm     | Parameters                               | Standard         |
|---------------------------|---------------|------------------------------------------|------------------|
| Password authentication   | OPAQUE        | RFC 9807, aPAKE                          | RFC 9807         |
| Password key derivation   | PBKDF2-SHA256 | 600,000 iterations, 32B salt, 32B output | NIST SP 800-132  |
| Vault encryption          | AES-256-GCM   | 256-bit key, 96-bit nonce, 128-bit tag   | NIST SP 800-38D  |
| Classical key exchange    | X25519        | Curve25519, 256-bit                      | RFC 7748         |
| Post-quantum KEM          | ML-KEM-1024   | Kyber-1024, NIST Level 5                 | NIST FIPS 203    |
| Hybrid key derivation     | HKDF-SHA256   | 64B IKM, info="trashmail-hybrid-kem-v1"  | RFC 5869         |
| Secret sharing            | Shamir SSS    | k=2, n=3, GF(2 <sup>8</sup> )            | Shamir (1979)    |
| Recovery key encoding     | BIP39         | 256-bit entropy, 24 words                | BIP-0039         |
| Recovery key derivation   | HKDF-SHA3-256 | Account-bound salt                       | NIST FIPS 202    |
| Recovery key verification | SHA3-256      | 32-byte output                           | NIST FIPS 202    |
| OPAQUE record encryption  | AES-256-GCM   | Server-side at-rest encryption           | NIST SP 800-38D  |
| Passkey vault unlock      | WebAuthn PRF  | HMAC-based, hardware-bound               | WebAuthn Level 2 |
| Compression               | gzip          | Level 6                                  | RFC 1952         |

| Component                     | Algorithm                    | Parameters                                                     | Standard                                       |
|-------------------------------|------------------------------|----------------------------------------------------------------|------------------------------------------------|
| Bucket padding                | Custom                       | 17 sizes (256B–16MB),<br>0xDEAD magic                          | —                                              |
| Response signing              | Ed25519                      | 256-bit key, 512-bit<br>signature                              | RFC 8032                                       |
| Archive hash chain            | SHA3-256                     | Per-block hash,<br>sequential linking                          | NIST FIPS 202                                  |
| Archive key wrapping<br>(CAK) | HKDF-SHA256 +<br>AES-256-GCM | Password-derived<br>wrapping key                               | RFC 5869 / NIST<br>SP 800-38D                  |
| Certificate<br>verification   | SHA-256                      | TLS cert fingerprint<br>comparison                             | —                                              |
| Email threading               | SHA-256                      | Hash of Message-ID                                             | NIST FIPS 180-4                                |
| Video-call group<br>keys      | MLS + X-Wing                 | Suite<br>AIONDA_PQ_CALL_V1<br>(0xFF01), ML-KEM-768<br>+ X25519 | RFC 9420 /<br>draft-conolly-cfrg-<br>xwing-kem |
| Video-call signatures         | ML-DSA-65                    | Dilithium, NIST Level 3                                        | NIST FIPS 204                                  |
| Video-call media              | SFrame<br>(AES-256-GCM)      | Per-frame, key from<br>MLS exporter                            | RFC 9605                                       |

## 20.2 Security Levels

| Algorithm                | Classical Security     | Post-Quantum Security             |
|--------------------------|------------------------|-----------------------------------|
| X25519                   | 128-bit                | Broken by Shor's algorithm        |
| ML-KEM-1024              | 256-bit equivalent     | NIST Level 5 ( $\approx$ AES-256) |
| AES-256-GCM              | 256-bit                | 128-bit (Grover's algorithm)      |
| SHA-256                  | 256-bit                | 128-bit (Grover's algorithm)      |
| SHA3-256                 | 256-bit                | 128-bit (Grover's algorithm)      |
| Hybrid KEM<br>(combined) | 128-bit (X25519 bound) | Level 5 (ML-KEM bound)            |

## 21. Comparison with Other Providers

| Feature                  | Aionda Mail                                             | Tuta Mail                                       | Proton Mail    |
|--------------------------|---------------------------------------------------------|-------------------------------------------------|----------------|
| <b>Country</b>           | Germany (Stuttgart)                                     | Germany (Hannover)                              | Switzerland    |
| <b>Zero-Knowledge</b>    | Yes (OPAQUE +<br>client-side crypto)                    | Yes                                             | Yes            |
| <b>Post-Quantum</b>      | Yes (ML-KEM-1024 +<br>X25519 Hybrid)                    | Yes (Kyber-based)                               | In development |
| <b>Password protocol</b> | OPAQUE (RFC 9807)<br>— password never<br>leaves browser | bcrypt (password<br>sent to server over<br>TLS) | SRP-based      |
| <b>Subject encrypted</b> | Yes                                                     | Yes                                             | No             |

| Feature                              | Aionda Mail                            | Tuta Mail         | Proton Mail           |
|--------------------------------------|----------------------------------------|-------------------|-----------------------|
| <b>Headers encrypted</b>             | Yes                                    | Partial           | No                    |
| <b>Contact names encrypted</b>       | Yes (in vault)                         | Yes               | No                    |
| <b>Disposable email addresses</b>    | Yes (core feature, unlimited for Plus) | No                | Yes (via SimpleLogin) |
| <b>Browser addon</b>                 | Yes (Chrome + Firefox)                 | No                | Via SimpleLogin       |
| <b>Folder sharing</b>                | Yes (Hybrid KEM per recipient)         | Limited           | Yes                   |
| <b>Open source client</b>            | No                                     | Yes               | Yes                   |
| <b>Security audit</b>                | Planned                                | Yes               | Yes                   |
| <b>Password recovery</b>             | No (by design)                         | No (by design)    | No (by design)        |
| <b>Passkey support</b>               | Yes (FIDO2 + PRF)                      | Yes               | Yes                   |
| <b>PGP support</b>                   | Yes (incoming + outgoing)              | No (own protocol) | Yes (OpenPGP)         |
| <b>GoBD-compliant email archive</b>  | Yes (SHA3-256 hash chain + Hybrid KEM) | No                | No                    |
| <b>MITM detection (browser ext.)</b> | Yes (Ed25519 signatures + TLS check)   | No                | No                    |
| <b>Perfect Forward Secrecy (API)</b> | Yes (per-request ephemeral keys)       | Unknown           | Unknown               |
| <b>Email size obfuscation</b>        | Yes (bucket padding)                   | Unknown           | No                    |

## 22. Limitations & Honest Boundaries

### 22.1 Web Application Trust Model

Aionda Mail is a web application. On every page load, the browser downloads JavaScript from our servers. A sophisticated attacker who compromises our servers could theoretically serve modified JavaScript that exfiltrates keys.

#### Current mitigations:

- Subresource Integrity (SRI) hashes on all script tags
- Content Security Policy (CSP) headers restrict script sources
- All critical cryptographic code is included in the main application bundle
- **Guardian browser extension** (Section 17): Ed25519 signature verification on all API responses detects server-side tampering; TLS certificate verification (Firefox) detects MITM proxies

**Planned mitigations:**

- Service Worker caching for offline operation (reduces trust-on-load frequency)

**22.2 Metadata Visibility**

While email content is fully encrypted, certain metadata is visible to the server:

- When emails were received (timestamps)
- Which DEA address received the email
- Approximate email size (within bucket boundaries)
- Account activity patterns

**22.3 Email Processing Log**

For diagnostic purposes, Aionda Mail includes an optional **email processing log** that can temporarily store the raw content of incoming emails. This feature is configurable per disposable email address (DEA) and can be enabled or disabled in the DEA settings (“Log email content”).

When enabled (opt-in per DEA):

- The full raw SMTP message (headers + body) is stored in plaintext on the server
- Automatic deletion after a short retention period (less than 7 days)
- Accessible only to the account owner via authenticated API
- Purpose: troubleshooting delivery issues, verifying forwarding, reviewing spam filtering decisions

When disabled:

- No email content is stored in the processing log
- Only metadata is logged (sender address, timestamp, delivery status)
- Vault encryption remains the sole storage mechanism

**Important:** This processing log is independent of the encrypted vault. Emails stored in the vault are always encrypted with Hybrid KEM regardless of the log setting. The processing log exists as a legacy feature from the email forwarding system and provides operational transparency. Users who require strict zero-knowledge storage for all emails should disable this option.

**22.4 External Email Security**

Emails sent to or received from non-Aionda addresses travel through the standard email infrastructure (SMTP). While stored encrypted in the vault, the email content was visible during transit unless PGP encryption was used.

**22.5 No Key Escrow**

There is no master key, backdoor, or recovery mechanism available to Aionda GmbH. If a user loses their password and all recovery methods, their data is permanently lost. This is an intentional design decision that proves the integrity of the zero-knowledge model.

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**23. Roadmap**

| Milestone                                 | Status    | Target |
|-------------------------------------------|-----------|--------|
| Zero-Knowledge Architecture               | Completed | —      |
| Post-Quantum Hybrid KEM (ML-KEM-1024)     | Completed | —      |
| OPAQUE Authentication (RFC 9807)          | Completed | —      |
| Shamir Secret Sharing (2-of-3)            | Completed | —      |
| Encrypted API Transport Layer             | Completed | —      |
| Passkey/WebAuthn PRF Support              | Completed | —      |
| End-to-End Encrypted Calendar             | Completed | —      |
| GoBD-Compliant Email Archive (Blockchain) | Completed | —      |
| Guardian MITM Protection (Ed25519)        | Completed | —      |
| TLS Certificate Verification (Firefox)    | Completed | —      |

## Document History

| Version | Date       | Changes                                                                                                                                                                                                                                                                                                                                                                                                                                                                                   |
|---------|------------|-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|
| 1.0     | March 2026 | Initial publication                                                                                                                                                                                                                                                                                                                                                                                                                                                                       |
| 1.1     | April 2026 | New chapter 10: Vault Drive (per-file-key architecture, sharing via key rewrap)                                                                                                                                                                                                                                                                                                                                                                                                           |
| 1.2     | April 2026 | Section 10.11: External Sharing — public share links with hybrid KEM envelopes, unlock_token flow, Argon2id password mode, URL fragment for link-only, tamper-evident audit chain (EXT_SHARE_*), and user-chosen link distribution                                                                                                                                                                                                                                                        |
| 1.3     | May 2026   | New chapter 11: Aionda Chat — post-quantum E2EE real-time messaging via AAR (Aionda Async Ratchet), a hybrid X25519 + ML-KEM-1024 variant of Signal X3DH + Double Ratchet. Per-message forward secrecy, post-compromise security, and integration with the Enterprise audit chain. Subsequent chapters renumbered 12-22. PT-BR and PT-PT translations added. Fixed double-numbering bug in PDF rendering (LaTeX auto-numbering disabled; chapter numbers now sourced from headings only). |

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| Version | Date      | Changes                                                                                                                                                                                                                                                                                                                                                                              |
|---------|-----------|--------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|
| 1.4     | June 2026 | New chapter 12: Aionda Video Calls — post-quantum E2EE group conferencing. MLS (RFC 9420) group key agreement with the AIONDA_PQ_CALL_V1 cipher suite (X-Wing = ML-KEM-768 + X25519, ML-DSA-65 signatures), SFrame (RFC 9605) per-frame AES-256-GCM media encryption keyed from the MLS exporter, SAS binding, and a zero-knowledge SFU relay. Subsequent chapters renumbered 13-23. |

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## Contact

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*This document describes the security architecture of Aionda Mail as of June 2026. Cryptographic systems evolve — this document will be updated as the architecture changes.*